

### **IOActive Security Advisory**

Title	
Severity	Critical
Discovered by	Mike Davis
Advisory Date	April 30, 2015

#### **Affected Products**

1. CyberLock CyberKey based access control solutions.

#### Overview

CyberLock offers a line of "high security" locks and cylinders as well as related products and services for updating, managing, provisioning, and storing CyberKeys. In various marketing materials, CyberKey is described as "unclonable" and suitable for use in money handling and critical infrastructure systems as a secure and auditable solution.

However, after some reverse engineering it appears that these devices are easily cloned, and new keys can be created from lost cylinders and keys regardless of the permissions granted to the key. Additionally, time-of-day restrictions are enforced by the key, not the cylinder, allowing an attacker access at any time regardless of the configuration.

#### Issues

- 1. By intercepting communications between any previously authorized CyberKey and any CyberLock, the site key can be extracted from the lock and used to create cloned keys.
- 2. All profile-based restrictions (time of day, one time access) are based entirely on the logic of the key itself and can be modified by an attacker who has produced a clone key.
- 3. While the audit trail stored by the lock provides some information in the event of unauthorized access, the lock itself is not aware of which key IDs are valid. Thus, an attacker can fill the log with nonsensical accesses.
- 4. Site keys can be recovered from cylinders and are stored in cleartext.
- 5. The "encryption" (encoding) algorithm used does not sufficiently protect credentials or enforce authenticity.
- 6. CyberLock cylinders can be relatively easily removed from CyberLock-branded padlock enclosures. With a few sharp strikes to the mechanical retainer, it will shear off, allowing the cylinder to be removed and shackle unlatched.



7. While the issue has already been reported, we feel these issues are exacerbated due to the already known "magnet" bypass (https://www.youtube.com/watch?v=YfldDq48I9U).

#### **Encryption Algorithm**

CODE:01E9 decrypt_super:			;	CODE XREF:
CODE:0133^Xp				
CODE:01E9	movlw	(byte_DATA_58)		
CODE:01EA	movwf	BANK0:FSR		
CODE:01EB	movlw	8		
CODE:01EC	movwf	arg_0		
CODE:01ED	swapf	byte_DATA_71, w		
CODE:01EE	addwf	byte_DATA_71, w		
CODE:01EF	subwf	byte_DATA_72, w		
CODE:01F0	addwf	byte DATA 73, w		
CODE:01F1	addwf	byte DATA 74, w		
CODE:01F2	subwf	byte DATA 75, w		
CODE:01F3	subwf	byte DATA 76, w		
CODE:01F4	movwf	byte DATA 26		
CODE:01F5 loc CODE 1F5:			;	CODE XREF:
decrypt super+14j				
CODE:01F5	movfw	BANKO INDF		
CODE:01F6	movwf	byte DATA 25	;	store the
byte				
CODE:01F7	movfw	byte DATA 26	;	get our key
ready				
CODE:01F8	subwf	BANKO INDF, f	;	subtract our
key from the byte		—		
CODE:01F9	movfw	byte DATA 25	;	get our now
modified byte back into w				
CODE:01FA	movwf	byte DATA 26	;	store our
newly calculated byte as t	the key f	for the next byte	Э	
CODE:01FB	incf	BANK0:FSR, f	;	move on to
the next byte				
CODE:01FC	decfsz	arg O, f		
CODE:01FD	b	loc CODE 1F5		

#### **Technical Details**

The CyberLock uses a standard memcmp-style key comparison function, which internally checks the 64-bit site key after "decryption." As the speed of the response to an invalid key (error code 0x10) is primarily dependent on the number of characters correctly matched, it may be possible to extract the site key through a brute-force attack on a CyberLock in-situ. In (limited) practice, however, it appears that accesses to flash memory adds significant noise to this timing signal and an attack would need to gather enough samples to average these signals out. Additionally, the CyberLock goes into an infinite loop, and the lock itself must be de-powered between attempts, somewhat mitigating this issue.

As the CyberLock is directly powered through the communications port, it appears that an SPA (power analysis) attack may succeed against a CyberLock in-situ, as the lock leaks a significant power side-channel to any potential "key" as the processor slowly clocks the key across an I<sup>2</sup>C bus at the Fcpu/4 bps.

However, this approach seems somewhat overboard given the existing issues.

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Figure 1: Power draw of initial flash access between marker A1 and A2 during key load (A2 marks the start of the lock ID message 0x02)



Figure 2: Disassembled CyberLock

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Figure 3: CyberKey firmware extraction methodology



Figure 4: CyberLock fuse clearing methodology for firmware extraction



#### **Responsible Disclosure**

- Research period: January 17, 2015 to March 15, 2015
- Vulnerability discovered: March 15, 2015
- Findings finalized: March 30, 2015
- First notification: March 31, 2015 to Bruce Stephenson, Senior Security Engineer in R&D
- Second notification: April 1, 2015 to <u>Support@cyberlock.com</u>
- Third notification: April 9, 2015 to CyberLock sales
- Fourth notification: April 11, 2015 to Tammy (media relations contact) *Email delivery confirmation received*.
- Fifth notification: April 17, 2015 to CyberLock sales and support
- Sixth notification: April 19, 2015 to CyberLock support